

Binarity, branchingness, and size effects

Jennifer Bellik (jbellik@ucsc.edu) & Nick Kalivoda (nkalivod@ucsc.edu)



Introduction

- In prosody, longer strings tend to be parsed into more constituents.
 - Ex. two feet in (abra)ka(dabra) but one foot in sha(zam)
- Such *size effects* are commonly captured using BINARITY constraints (e.g., Inkelas & Zec 1990, Ito & Mester 1992, Prince & Smolensky 1993/2004; Sandalo & Truckenbrodt 2002, Prieto 2007, Selkirk 2011, Elfner 2012).
- Implementations of BINARITY come in two major flavors, which often aren't distinguished:
 - Branch-counting BINARITY** (formalized in [1]) requires a node to branch into two children (of any category), and echoes syntactic notions of binarity.
 - Leaf-counting BINARITY** (formalized in [2]) reflects a rhythmic conception of binarity, derived from foot-building. It requires a node to contain two dominated nodes of some particular category (e.g. σ), and sometimes lower categories.

- BINARITY-BRANCHES(K) = BIN-BR(K): Assign a violation for every node of category K with more than two branches (immediate children, of any category).
- BINARITY-LEAVES(K,L) = BIN-LV(K, L): Assign a violation for every node of category K that dominates more than two nodes of category L at any level, where $L < K$.

Big question: Are both versions of BINARITY necessary and desirable for understanding the syntax-prosody interface?

Proposal:

- Binarity in prosody is best conceived of as counting branches, regardless of category.
- Category-sensitive leaf-counting binarity should be restricted to counting rhythmic categories.

Binarity under Weak vs. Strict Layering

Under Strict Layering: No difference between branch- and leaf-counting

- In structures that conform to *Strict Layering* (Selkirk 1984), no level-skipping (non-exhaustive parsing) or level-doubling (recursion) is permitted.
- Therefore, every branch in a Strictly Layered tree corresponds to a child of the next lower prosodic category, and counting branches is equivalent to counting leaves of the next lower prosodic category.
 - Ex. Strictly Layered (3a,d) incur the same violations under branch- and leaf-counting binarity.
- So in effect, branch-counting binarity is already employed in analyses with Strict Layering, such as Shih (2017), Prieto (2007), and Sandalo & Truckenbrodt (2002), even when the analysis defines binarity in leaf-counting terms.

Under Weak Layering: Branch-counting and leaf-counting pull apart

- Weak Layering* (Ito & Mester 1992/2003) permits recursion (3c) and level skipping (3b)
- As a result, not every branch corresponds to a child of the next lower prosodic category.
 - Recursion produces children of the same category as the parent
 - Non-exhaustive parsing produces children of an even lower prosodic category
 - Ex.: in (3c), ϕ_1 has only two branches (to ϕ_2 and ϕ_3) and satisfies BIN-BR, but ϕ_1 violates BIN-LV since it dominates four leaves ($\omega_1, \omega_2, \omega_3, \omega_4$).

(3) Constraint comparison

	$[[XP X_1 [XP X_2 [XP X_3 X_4]]]]$	BIN(BR)	BIN(LV)	Layering
a.				strict
b.				weak (+EXH)
c.			ϕ_1	weak (+NONREC)
d.			ϕ	strict

Selected References. Bellik, J. & N. Kalivoda. 2016. Adjunction and Branchingness Effects in Syntax-Prosody Mapping. In Hansson, G.O., A Farris-Trimble, K. McMullin, & D. Pulleyblank (eds.), *Supplemental Proceedings of the 2015 Annual Meeting on Phonology*. Elfner, E. 2012. *Syntax-Prosody Interactions in Irish*. University of Massachusetts - Amherst dissertation. Ito, J. & A. Mester. 1992/2003. *Weak Layering and Word Binarity*. Linguistic Research Center, LRC-92-09, University of California, Santa Cruz. — 2013. Prosodic subcategories in Japanese. *Lingua* 124, pp. 20-40. — 2015. The perfect prosodic word in Danish. *Nordic Journal of Linguistics*, Vol. 38, No. 1, pp. 5-36. Kalivoda, N. 2018. *Syntax-Prosody Mismatches in Optimality Theory*. Ph.D. Thesis, UC Santa Cruz. Prince, A. & P. Smolensky. 1993/2004. *Optimality Theory: Constraint Interaction in Generative Grammar*. Selkirk, E. 1984. *Phonology and Syntax: The Relation between Sound and Structure*. MIT Press. — 2011. The syntax-phonology interface. In J. Goldsmith, J. Riggle & A. Yu (eds.) *The Handbook of Phonological Theory*, 2nd edition.

Branch-counting motivates size-driven recursion

Branch-counting binarity (1) assigns a violation to nodes that branch into more than two immediate children. Therefore, it prefers candidates with more prosodic structure.

Case study 1: Size-effects in Danish compound words

- Danish glottal accent ('stød') diagnoses the right edge of a prosodic word in Danish
 - Stød reveals length-driven differences in compound phrasing (Ito & Mester 2015):
- (4) a. $[_\omega \text{ to:g } [_\omega \text{ passa:ge:r}]]$ 'train passenger' b. $[_\omega [_\omega \text{ passa:ge:r}]]$ 'pass. train'
- These size effects can be derived with the ranking BIN-BR \gg NONREC \gg MATCH (5).

(5) Tableau for Danish 'train passenger'

	$[[XP [XP [XP X_1 X_2]]]]$	BIN-BR	NONREC	MATCH	BIN(ω, Ft)
a. \rightarrow			*	*	*
b.		**!			*
c.		*!		**	*

Case study 2: Size-effects in Irish phrasing

In Irish, LH shows the left boundary of ϕ_{NonMin} and HL shows the right boundary of ϕ (from Elfner 2012). This phrasing is governed by the constraint ranking in (6) (Elfner 2012).

- $[[XP V [TP [NP N] [VP [NP NA]]]] \rightarrow (\phi (VN) (\phi NA))$
- Elfner attributes this to **STRONGSTART** (Selkirk 2011), which penalizes $(\phi \omega \phi \dots)$ structures, and outranks Match.
- But **STRONGSTART** violations are tolerated in order to avoid **BIN-BR** violations (6).
- To avoid a binarity violation in winner (6a), binarity must be assessed by counting branches for ϕ_1 .
- (6a)'s ϕ_1 , which dominates five ω s, still violates leaf-counting **BIN- ω** , but is binary branching.

(6) Tableau for Irish (Elfner 2012)

	$[[XP V [TP [NP N_1 As] [VP [NP N_2 Ao]]]]]$	BIN-BR	STST	MATCH	BIN- ω
a. \rightarrow			ϕ_1		ϕ_1, ϕ_2
b.			ϕ_1, ϕ_2	TP, DP _S	ϕ_1, ϕ_2
c.		$\phi_2!$		TP, DP _S	ϕ_1, ϕ_2

Similar examples are found in Kimatuumbi (Kalivoda 2018) and Mandarin (Shih 2017).

Upshot: Only branch-counting motivates size-effects and a closer syntax-prosody match.

Leaf-counting motivates size-driven category change

Leaf-counting binarity counts dominated nodes of a lower prosodic category (cf. Dresher & van der Hulst 1998), and can motivate a size-driven change of category.

- In (3a,b), the violation of leaf-counting BIN-LV(ϕ, ω) is avoided by punting the binarity violation up to the level of ι , so that (3a,b) outperform their counterparts (3c,d).
- Unlike BIN-LV, BIN-BR does not favor the level-skipping (3b) over (3a,c)—it doesn't motivate a category change for the root prosodic node.
- The case study shown below, where Ft- and σ -counting BIN-LV promotes ω to ϕ , is the only instance of such a category change known to us.

Case study: Japanese compound phrasing

Category-promotion can be seen in Japanese compounds (Ito & Mester 2013).

- Compounds consisting of no more than two feet are parsed into a prosodic structure rooted in ω (as diagnosed by ω -internal rendaku voicing, compound accent, etc.)
- But in compounds of the form $[[Ft][\sigma Ft]]$, the category of the root node changes to ϕ instead (7a), so that no ω dominates more than two feet or syllables.
- This category-change is driven by leaf-counting binarity (BIN-LV($\omega, [Ft, \sigma]$)): each ω can dominate only up to two nodes of lower prosodic categories.
- Branch-counting binarity can't drive this category-change—both (7a,b) satisfy BIN-BR.

(7) Tableau for Japanese compound

	$[[Ft][\sigma Ft]]$	BIN ($\omega, [Ft, \sigma]$)	BIN (ω, BR)
a. \rightarrow			
b.		*!	

If category-promotion when only occurs when rhythmic categories, not interface categories, are counted, then CON should only include versions of BIN-LV that count feet and syllables, not ω or ϕ (i.e., BIN-LV($\omega, [Ft, \sigma]$), but not BIN-LV(ϕ, ω) or BIN-LV(ι, ϕ))

Leaf-counting's undesirable typological predictions

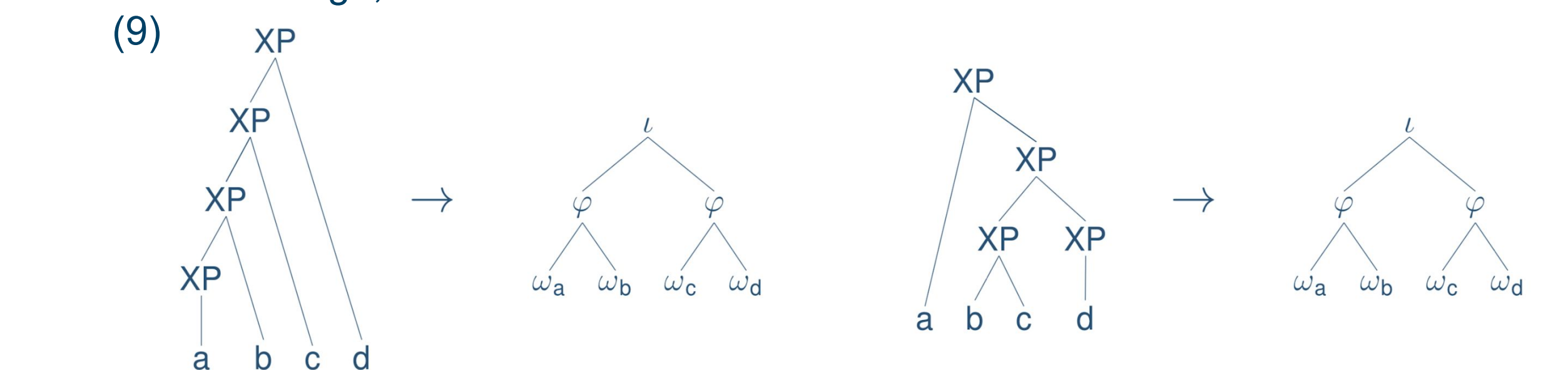
Typologies are larger under leaf-counting binarity. In a case study of Kinyambo phrasing, leaf-counting systems predicted typologies that were, on average, 80% larger than branch-counting typologies, with an average of 10.67 more languages (Bellik & Kalivoda 2016).

(8) Number of languages in Bellik & Kalivoda (2016) typologies (4 inputs from Kinyambo)

	Pair 1 (Match, low)	Pair 2 (Align, low)	Pair 3 (Match, high)	Pair 4 (Align, high)	Pair 5 (Match, high & low)	Pair 6 (Match, high & low)
BIN-LV lgs	19	40	17	30	15	28
BIN-BR lgs	17	15	18	12	11	12
$[[BR]-[LV]]$ ($[[BR]-[LV]]/ BR $)	2 (12%)	25 (167%)	-1 (-5%)	18 (150%)	4 (36%)	16 (133%)

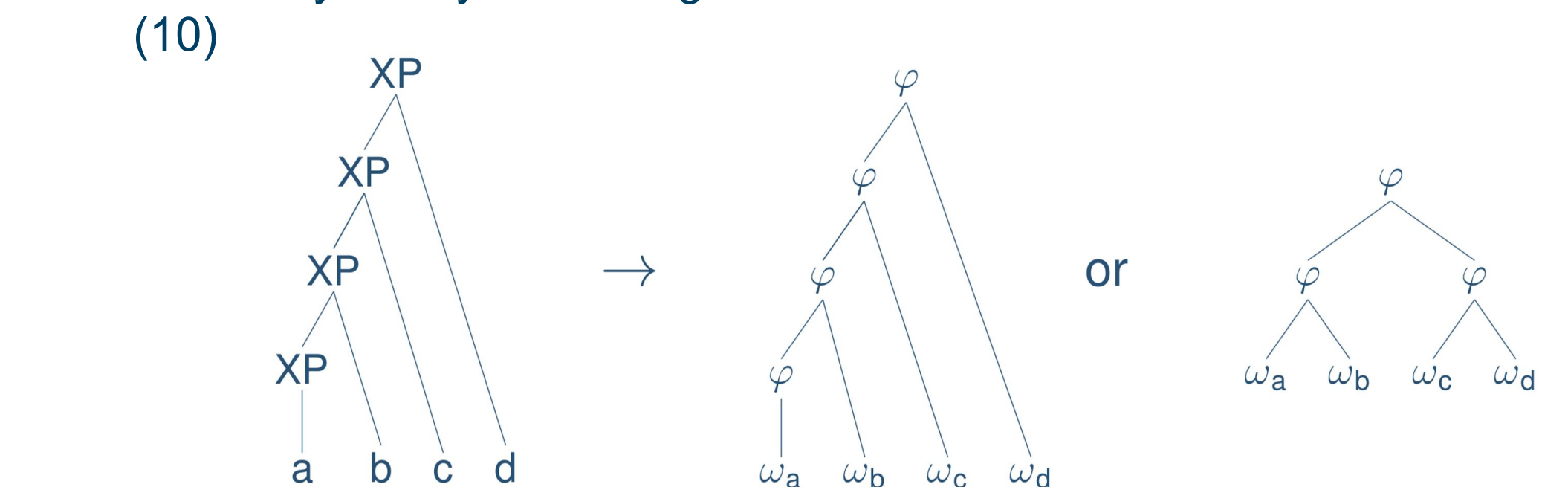
Leaf-counting binarity predicts a language that completely disregards syntax.

- If BIN-LV were top-ranked, syntaxes like $[a [b [c [d]]]]$, $[[ab][cd]]$ and $[[[[a] b] c] d]$ would all be optimally parsed as $(\iota (\phi a b) (\phi c d))$. Note the mapping of the top XP to ι , not ϕ . To our knowledge, this is unattested.



In Japanese, $[[[[a] b] c] d] \rightarrow (\iota (\phi a b) (\phi c d))$ (Kubozono 1989). However, this is not the same as $(\iota (\phi ab) (\phi cd))$ above, since there is a ϕ containing a, b, c, d. This 4-leaved ϕ violates BIN-LV, so this is not an argument in favor of BIN-LV *per se* (Kalivoda 2018).

- In contrast, BIN-BR never compels syntax-prosody mismatches since the syntactic input is already binary-branching.



Further disadvantages of leaf-counting binarity

Leaf-counting binarity is usually redundant, even under Weak Layering

- Leaf-counting is only essential if it rules out an intended loser that does *not* violate branch-counting. Call such phrasings that violate BIN-LV but not BIN-BR **leaf-violators**.
- Most leaf-violators can be ruled out by other prosodic well-formedness constraints. For all phrasings of three words (generated using SPOT [Bellik, Bellik & Kalivoda 2018]), *all* leaf-violators incurred additional penalties from **STRONGSTART**, **EQUALSISTERS**, or **NONRECURSIVITY**, compared to all non-leaf-violating possible optima.

Leaf-counting is more computationally complex. Branch-counting only examines a node's immediate children—a local search. But leaf-counting requires a global search through theoretically unbounded levels of recursion for all nodes of some lower category.

Conclusion

Branch-counting and leaf-counting versions of binarity differ significantly in their predictions:

- Size effects:** Branch-counting motivates size-driven recursion or splitting, deriving the desired size effects. Leaf-counting cannot predict such size effects.
- Category change:** Leaf-counting does motivate size-driven category change, which is only attested for ω being promoted to ϕ .
- Typology:** Branch-counting predicts smaller typologies than leaf-counting, and leaf-counting predicts a language we believe to be unattested.

Furthermore, leaf-counting is almost always redundant. CON and the size of predicted typologies could be constrained by restricting BIN-LV to counting rhythmic categories (σ, Ft), while allowing branch-counting binarity to apply at all levels of the prosodic hierarchy.